

## Lecture # 2, Quantum Computation 2: QEC Criteria

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Outline:

0. Review
1. Classical Coding
2. Q. Coding
3. Operator Measurement and Error Syndromes
4. Shor 9 Qubit Code
5. Quantum Error correction Codes Criteria (QEC criteria)

## 0. Review

$$\rho \xrightarrow{\varepsilon} \rho'$$
$$\varepsilon(\rho) = \sum_k E_k \rho E_k^\dagger \text{ where } \sum_k E_k E_k^\dagger = I$$

## 1. CLASSICAL CODING

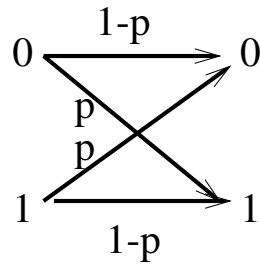


FIG. 1: a binary symmetric channel

$P$  = prob of error

Definition: A Classical  $[n,k,d]$  code is a set of  $2^k$   $n$ -bit strings which have a minimum Hamming distance  $d$ .

Definition: A Hamming distance between two bit strings is  $d(x, y) = w(x \oplus y)$  where  $\oplus$  is the x-or operator, and  $w$  is an operation that counts the number of ones.

Example:  $0_{L(\text{ogical})} = 000$ ,  $1_L = 111$  is a  $[3,1,3]$  code

could send	could receive	prob	decode	prob. of error
$0_L = 000$	000	$(1-p)^3$	0	
	001	$p(1-p)^2$	0	
	010	$p(1-p)^2$	0	
	100	$p(1-p)^2$	0	
	011	$p^2(1-p)$	1	$p^2(1-p)+$
	101	$p^2(1-p)$	1	$p^2(1-p)+$
	110	$p^2(1-p)$	1	$p^2(1-p)+$
	111	$p^3$	1	$p^3$

So, the total probability of error is  $3p^2 - 2p^3 = O(p^2)$

## 2. QUANTUM CODING

1995: Thought error correction to be impossible!

1. States collapse on measurement
2. Classically error occurs or does not occur. In Q. M., errors are continuous:  $\alpha|0\rangle + \beta|1\rangle \rightarrow (\alpha + \varepsilon)|0\rangle + \dots$
3. No cloning Thm prohibits copying, so cannot create  $\alpha|0\rangle + \beta|1\rangle \rightarrow (\alpha|0\rangle + \beta|1\rangle)(\alpha|0\rangle + \beta|1\rangle)(\alpha|0\rangle + \beta|1\rangle)$

The Solutions:

1. Measure only the effect of the environment, not the state (i.e. did an error occur?)
2. & 3. Orthogonalize errors using entanglement: the environment has done one thing, or another, in an entangled way.  $\alpha|didsomething\rangle + \beta|didnothing\rangle$

Example: The Quantum Bit Flip Code:

$$|0_L\rangle = |000\rangle$$

$$|1_L\rangle = |111\rangle$$

$$|\Psi_L\rangle = \alpha|0_L\rangle + \beta|1_L\rangle$$

suppose  $\varepsilon(\rho) = (1 - P)\rho + PX\rho X$  where  $P$  is the probability of error and  $X$  is the error operator.

$$A \xrightarrow{\varepsilon(\rho)} B$$

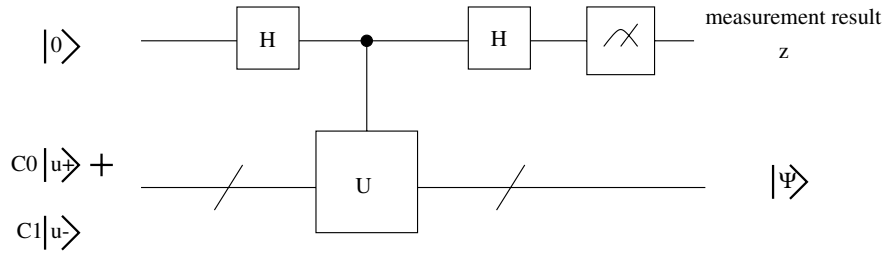
Define: An  $[[n,k]]$  quantum code  $C$  is a  $k$ -qubit subspace of an  $n$ -qubit Hilbert space. So, for our example,  $k=3$ ,  $n=1$ .

Input	$\xrightarrow{\varepsilon^{\otimes 3}}$ Output	prob	decode	prob. of error
$ \Psi\rangle = \alpha 000\rangle + \beta 111\rangle$	$\alpha 000\rangle + \beta 111\rangle$	$(1-p)^3$	0	
	$\alpha 001\rangle + \beta 110\rangle$	$p(1-p)^2$	0	
	$\alpha 010\rangle + \beta 101\rangle$	$p(1-p)^2$	0	
	$\alpha 100\rangle + \beta 011\rangle$	$p(1-p)^2$	0	
	$\alpha 011\rangle + \beta 100\rangle$	$p^2(1-p)$	1	$p^2(1-p)+$
	$\alpha 101\rangle + \beta 010\rangle$	$p^2(1-p)$	1	$p^2(1-p)+$
	$\alpha 110\rangle + \beta 001\rangle$	$p^2(1-p)$	1	$p^2(1-p)+$
	$\alpha 111\rangle + \beta 000\rangle$	$p^3$	1	$p^3$

### 3. OPERATOR MEASUREMENT

Given  $U$  with eigenvalues  $\pm 1$ , eigenvectors  $|u_{\pm}\rangle$

Definition: Measuring  $U$



Initially, the state is  $|0\rangle(C_0|u_+\rangle + C_1|u_-\rangle)$

After the first Hadamard,  $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)(C_0|u_+\rangle + C_1|u_-\rangle)$

After the Controlled-U gate,  $\frac{1}{\sqrt{2}}(|0\rangle(C_0|u_+\rangle + C_1|u_-\rangle) + |1\rangle(C_0|u_+\rangle - C_1|u_-\rangle))$

After the last Hadamard,  $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)(C_0|u_+\rangle + C_1|u_-\rangle) + (|0\rangle - |1\rangle)(C_0|u_+\rangle - C_1|u_-\rangle)$   
 $= |0\rangle C_0|u_+\rangle + |1\rangle C_1|u_-\rangle$

If the measurement is  $z = 0$ , then  $|Psi\rangle = |u_+\rangle$ . (With prob  $C_0^2$ ,  $z = 0$ )

If the measurement is  $z = 1$ , then  $|Psi\rangle = |u_-\rangle$ . (With prob  $C_1^2$ ,  $z = 1$ )

#### 3.1. Error Correction Syndrome Measurement

$$U_1 = \sigma_z^1 \sigma_z^2 = \sigma_z \sigma_z I$$

$$U_2 = \sigma_z^2 \sigma_z^3 = I \sigma_z \sigma_z$$

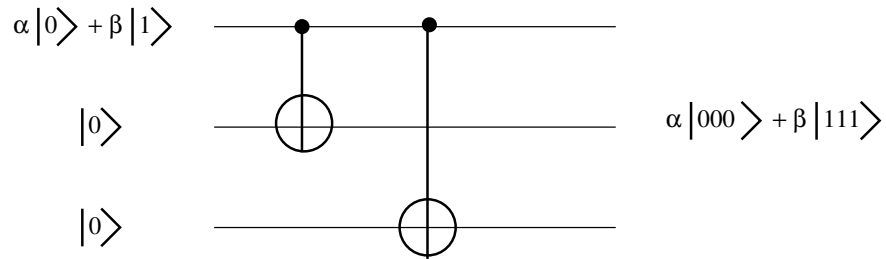
state	$U_1$	$U_2$
$\alpha 000\rangle + \beta 111\rangle$	0	0
$\alpha 001\rangle + \beta 110\rangle$	0	1
$\alpha 010\rangle + \beta 101\rangle$	1	1
$\alpha 100\rangle + \beta 011\rangle$	1	0

TABLE I: 0 represents a +1 eigenstate of  $U_i$ , and 1 represents a -1 eigenstate.

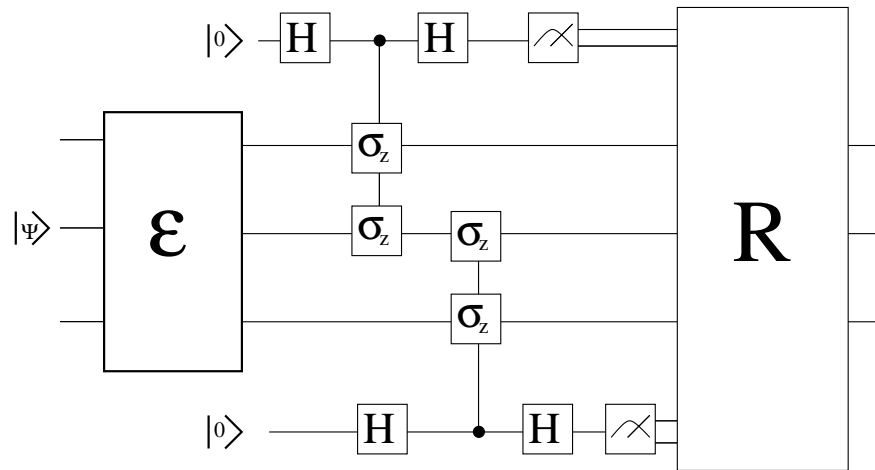
Steps to Error Correction:

1. measure syndrome operators (here,  $U_1$  &  $U_2$ )
2. Apply recovery operator R (here,  $00 \rightarrow I$ ,  $01 \rightarrow \sigma_x^3$ ,  $11 \rightarrow \sigma_x^2$ ,  $10 \rightarrow \sigma_x^1$ )

To create the initial state  $|\Psi_L\rangle$ :



And then to error correct:



note: the double lines indicate classical information.

Claim:

This scheme also corrects for a small continuous rotation error!

We will do this on one bit to demonstrate.

$$\varepsilon(\rho) = e^{i\epsilon\sigma_x}\rho e^{i\epsilon\sigma_x}$$

$$e^{-\epsilon\sigma_x} = R_x(2\epsilon)$$

$$R_{x^1}(2\epsilon)|\Psi\rangle \cong |\Psi\rangle - i\epsilon\sigma_x^1|Psi\rangle \equiv |\Psi'\rangle$$

$$\text{The fidelity is } F = \sqrt{|\langle\Psi|\Psi'\rangle|^2} \cong 1 - \epsilon$$

Syndrome measurement collapses error into either I or  $\sigma_x^1$

$$F(R(\varepsilon(\rho)), |\Psi\rangle) \cong? \cong 1 - \epsilon^2$$

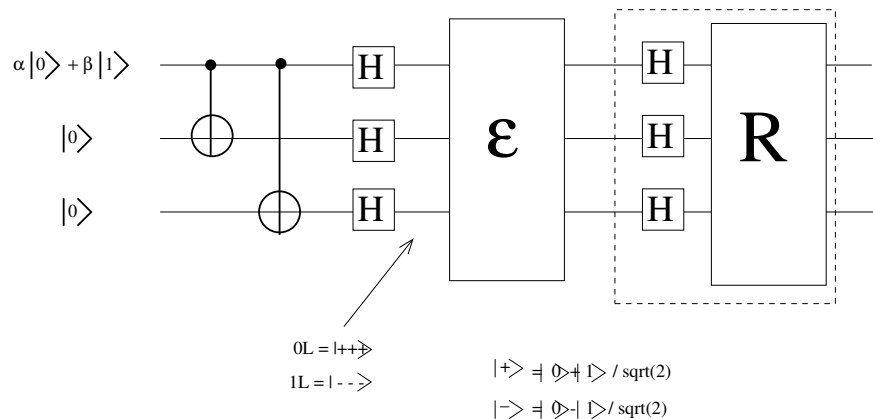
Example: The Phase Flip Code

$$\varepsilon_{\text{phaseflip}}(\rho) = (1 - P)\rho + P\sigma_z\rho\sigma_z$$

$$\text{Recall } H\sigma_xH = \sigma_z, H\sigma_zH = \sigma_x$$

$$\text{So, } H\varepsilon_{\text{phaseflip}}(H\rho H)H = \varepsilon_{\text{bitflip}}$$

Explicitly,



For the bit flip:  $U_o = \sigma_z\sigma_zI$  and  $U_1 = I\sigma_z\sigma_z$

For the phase flip:  $U_o = \sigma_x\sigma_xI$  and  $U_1 = I\sigma_x\sigma_x$

Claim:

arbitrary errors can be described as  $\sigma_x$ ,  $\sigma_z$ , and  $\sigma_x\sigma_z$  errors

Proof Argument:

$$\varepsilon(\rho) = \sum_k E_k\rho E_k^\dagger$$

where we are guaranteed  $\sum_k E_k E_k^\dagger = I$

Recall pauli matrices  $\sigma_j = I, \sigma_x, \sigma_y, \sigma_z$ , and that  $\sigma_y = -i\sigma_x\sigma_z$

Since  $\sigma_j$  is a basis for all 2x2 hermitian matrices, let  $E_k = \sum_j C_{kj}\sigma_j$ .

Then,  $\varepsilon(\rho) = \sum_{kjj'} C_{k_j} C_{k_{j'}}^* \sigma_j \rho \sigma_{j'}$

$\boxed{\varepsilon(\rho) \sum_{jj'} \chi_{jj'} \sigma_j \rho \sigma_{j'}}$  is the “Chi representation or OSR”

Example: recall

$$R_x(2\varepsilon)|\Psi\rangle \cong |Psi\rangle - i\varepsilon\sigma_x|Psi\rangle$$

$$\varepsilon(\rho) = \rho - i\varepsilon\sigma_x\rho - i\varepsilon\rho\sigma_x + \varepsilon^2\sigma_x\rho\sigma_x$$

The  $-i\varepsilon\sigma_x\rho - i\varepsilon\rho\sigma_x$  term disappears in the syndrome measurement, and the  $\rho + \varepsilon^2\sigma_x\rho\sigma_x$  term remains.

The result is that the syndrome measurement projects the environment into a definite error state.



#### 4. SHOR 9 QUBIT CODE

$$|0_L\rangle = (|000\rangle + |111\rangle)^{\otimes 3} / \sqrt{8}$$

$$|1_L\rangle = (|000\rangle - |111\rangle)^{\otimes 3} / \sqrt{8}$$

this code will correct ANY single qubit error.

Syndrome Measurements:

for a bit flip:  $\sigma_z^1 \sigma_z^2, \sigma_z^2 \sigma_z^3, \sigma_z^4 \sigma_z^5, \sigma_z^5 \sigma_z^6, \sigma_z^7 \sigma_z^8, \sigma_z^8 \sigma_z^9$

for a phase flip:  $\sigma_x^1 \sigma_x^2, \sigma_x^3 \sigma_x^4, \sigma_x^4 \sigma_x^5, \sigma_x^5 \sigma_x^6, \sigma_x^6 \sigma_x^7, \sigma_x^8 \sigma_x^9$ ,

## 5. QEC CRITERIA/CONDITIONS

Channel:  $E(\rho) = \sum_k E_k \rho E_k^\dagger$

Thm: Let  $C$  be a quantum Code defined by the orthonormal states  $\{ |\Psi_l\rangle \}$

$\exists$  a quantum recovery operation  $R$  correction  $\varepsilon$  on  $C$  iff:

1. Orthogonality:

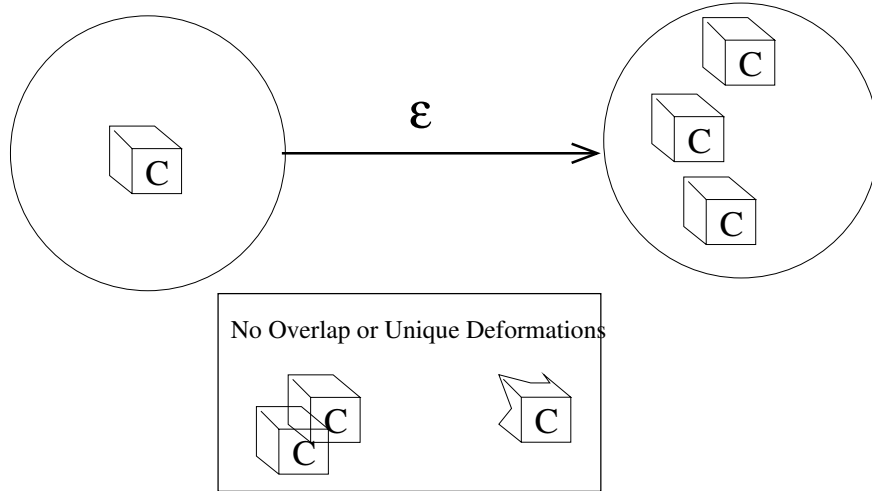
if I have 2 errors  $j$  and  $k$ ,

$$\langle \Psi_l | E_j^\dagger E_k | \Psi_l \rangle = 0$$

2. Nondeformation criteria:

$$\langle \Psi_l | E_k^\dagger E_k | \Psi_l \rangle = d_k \forall l$$

this is so you cannot distinguish shrinking on different code words, all shrinking is the same.



note that  $d_k$  implies probability loss, but not information loss,  $\sum_k d_k = 1$  since  $\sum_k E_k^\dagger E_k = 1$

Proof: ( $\rightarrow$ )

Let  $P = \sum_l |\Psi_l\rangle\langle\Psi_l|$  (project onto  $C$ )

note  $PE_j^\dagger E_k P = d_k \delta_{jk} P$  (\*)

note by Polar decomposition (extracting rotation and shrinkage)  $E_k P = U_k \sqrt{PE_k^\dagger E_k P} = \sqrt{d_k} U_k P$  where  $\sqrt{d_k}$  is the shrinkage and  $U_k P$  is the rotation.

1. Syndrome measurement:

$$\text{let } P_k = U_k P U_k^\dagger = \frac{E_k P U_k^\dagger}{\sqrt{d_k}} = \frac{U_k P E_k^\dagger}{\sqrt{d_k}}$$

By (\*), the  $P_k$ s are orthogonal:

$$\forall k \neq j, P_k P_j \propto U_k P E_k^\dagger E_j P U_j^\dagger = 0$$

measure  $P_k$  output  $k$  syndrome.

2. Apply Recovery R

$$R(\rho) = \sum_k U_k^\dagger P_k \rho P_k U_k$$

note for  $|\Psi\rangle \in C$ ,

$$U_k^\dagger P_k E_j |\Psi\rangle = \frac{U_k^\dagger U_k P E_k^\dagger E_j P |\Psi\rangle}{\sqrt{d_k}} = \frac{\delta_{jk} d_k P |\Psi\rangle}{\sqrt{d_k}} = \sqrt{d_k} \delta_{jk} |\Psi\rangle$$

Thus:

$$R(\varepsilon(|\Psi\rangle\langle\Psi|)) = R(\sum_j E_j |\Psi\rangle\langle\Psi| E_j^T) = \sum_{jk} U_k^\dagger P_k E_j^\dagger P_k U_k = \sum_{jk} d_k \delta_{jk} P = |\Psi\rangle\langle\Psi|$$