

## 18.335 Midterm, Fall 2012

### Problem 1: (25 points)

- (a) Your friend Alyssa P. Hacker claims that the function  $f(x) = \sin x$  can be computed accurately (small forward relative error) near  $x = 0$ , but not near  $x = 2\pi$ , despite the fact that the function is periodic in exact arithmetic. True or false? Why?
- (b) Matlab provides a function `log1p(x)` that computes  $\ln(1+x)$ . What is the point of providing such a function, as opposed to just letting the user compute  $\ln(1+x)$  herself? (Hint: not performance.) Outline a possible implementation of `log1p(x)` [rough pseudocode is fine].
- (c) Matlab provides a function `gamma(x)` that computes the “Gamma” function  $\Gamma(x) = \int_0^\infty e^{-t} t^{x-1} dt$ , which is a generalization of factorials, since  $\Gamma(n+1) = n!$ . Matlab also provides a function `gammaln(x)` that computes  $\ln[\Gamma(x)]$ . What is the point of providing a separate `gammaln` function? (Hint: not performance.)

### Problem 2: (5+10+10 points)

Recall that a floating-point implementation  $\tilde{f}(x)$  of a function  $f(x)$  (between two normed vector spaces) is said to be *backwards stable* if, for every  $x$ , there exists some  $\tilde{x}$  such that  $\tilde{f}(x) = f(\tilde{x})$  for  $\|\tilde{x} - x\| = \|x\|O(\epsilon_{\text{machine}})$ . Consider how you would apply this definition to a function  $f(x, y)$  of *two* arguments  $x$  and  $y$ . Two possibilities are:

- First: The most direct application of the original definition would be to define a single vector space on pairs  $v = (x, y)$  in the obvious way  $[(x_1, y_1) + (x_2, y_2) = (x_1 + x_2, y_1 + y_2)]$  and  $\alpha \cdot (x, y) = (\alpha x, \alpha y)$ , with some norm  $\|(x, y)\|$  on pairs. Then  $\tilde{f}$  is backwards stable if for every  $(x, y)$  there exist  $(\tilde{x}, \tilde{y})$  with  $\tilde{f}(x, y) = f(\tilde{x}, \tilde{y})$  and  $\|(\tilde{x}, \tilde{y}) - (x, y)\| = \|(x, y)\|O(\epsilon_{\text{machine}})$ .
  - Second: Alternatively, we could say  $\tilde{f}$  is backwards stable if for every  $x, y$  there exist  $\tilde{x}, \tilde{y}$  with  $\tilde{f}(x, y) = f(\tilde{x}, \tilde{y})$  and  $\|\tilde{x} - x\| = \|x\|O(\epsilon_{\text{machine}})$  and  $\|\tilde{y} - y\| = \|y\|O(\epsilon_{\text{machine}})$ .
- (a) Given norms  $\|x\|$  and  $\|y\|$  on  $x$  and  $y$ , give an example of a valid norm  $\|(x, y)\|$  on the vector space of pairs  $(x, y)$ .

- (b) Does First  $\implies$  Second, or Second  $\implies$  First, or both, or neither? Why?
- (c) In class, we proved that summation of  $n$  floating-point numbers, in some sequential order, is backwards stable. Suppose we sum  $m+n$  floating point numbers  $x \in \mathbb{R}^m$  and  $y \in \mathbb{R}^n$  by  $\tilde{f}(x, y) = x_1 \oplus x_2 \oplus x_3 \oplus \dots \oplus x_m \oplus y_1 \oplus y_2 \oplus \dots \oplus y_n$ , doing the floating-point additions ( $\oplus$ ) sequentially from left to right. Is this backwards stable in the First sense? In the Second sense? (No complicated proof required, but give a brief justification if true and a counterexample if false.)

### Problem 3: (25 points)

Say  $A$  is an  $m \times m$  diagonalizable matrix with eigenvectors  $x_1, x_2, \dots, x_m$  (normalized to  $\|x_k\|_2 = 1$  for convenience) and distinct-magnitude eigenvalues  $\lambda_k$  such that  $|\lambda_1| > |\lambda_2| > \dots > |\lambda_m|$ . In class, we showed that  $n$  steps of the QR algorithm produce a matrix  $A_n = Q^{(n)*} A Q^{(n)}$  where  $Q^{(n)}$  is equivalent (in exact arithmetic) to QR factorizing  $A^n = Q^{(n)} R^{(n)}$ . This proof was general for all  $A$ . For the specific case of  $A = A^*$  where the eigenvectors are orthonormal, we concluded that as  $n \rightarrow \infty$  we obtain  $Q^{(n)} \rightarrow$  eigenvectors  $(x_1 \dots x_m)$  and  $A_n \rightarrow \Lambda = \text{diag}(\lambda_1, \dots, \lambda_m)$ .

**Show** that if  $A \neq A^*$  (so that the eigenvectors  $x_k$  are no longer in generally orthogonal), the QR algorithm approaches  $A_n \rightarrow T$  and  $Q^{(n)} \rightarrow Q$  where  $T = Q^* A Q$  is the **Schur factorization** of  $A$ . (Hint: show that  $q_k = Q^{(n)} e_k$ , the  $k$ -th column of  $Q^{(n)}$ , is in the span  $\langle x_1, x_2, \dots, x_k \rangle$  as  $n \rightarrow \infty$ , by considering  $v_k = A^n e_k$ , the  $k$ -th column of  $A^n$ . Similar to class, think about the power method  $A^n e_k$ , and what Gram-Schmidt does to this.)

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