15.081J/6.251J Introduction to Mathematical Programming

Lecture 7: The Simplex Method III

1 Outline

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- Finding an initial BFS
- The complete algorithm
- The column geometry
- Computational efficiency
- The diameter of polyhedra and the Hirch conjecture

2 Finding an initial BFS

SLIDE 2

- **Goal:** Obtain a BFS of $\boldsymbol{Ax} = \boldsymbol{b}, \quad \boldsymbol{x} \geq \boldsymbol{0}$
  or decide that LOP is infeasible.
- Special case: $\boldsymbol{b} \geq \boldsymbol{0}$

$$\boldsymbol{Ax} \leq \boldsymbol{b}, \quad \boldsymbol{x} \geq \boldsymbol{0}$$

$$\Rightarrow \boldsymbol{Ax} + \boldsymbol{s} = \boldsymbol{b}, \quad \boldsymbol{x}, \boldsymbol{s} \geq \boldsymbol{0}$$

$$\boldsymbol{s} = \boldsymbol{b}, \quad \boldsymbol{x} = \boldsymbol{0}$$

2.1 Artificial variables

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$$\boldsymbol{Ax} = \boldsymbol{b}, \quad \boldsymbol{x} \geq \boldsymbol{0}$$

1. Multiply rows with -1 to get $\boldsymbol{b} \geq \boldsymbol{0}$.
2. Introduce artificial variables $\boldsymbol{y}$, start with initial BFS $\boldsymbol{y} = \boldsymbol{b}, \ \boldsymbol{x} = \boldsymbol{0}$, and apply simplex to auxiliary problem

$$\begin{array}{ll} \min & y_1 + y_2 + \ldots + y_m \\ \text{s.t.} & \boldsymbol{Ax} + \boldsymbol{y} = \boldsymbol{b} \\ & \boldsymbol{x}, \boldsymbol{y} \geq \boldsymbol{0} \end{array}$$

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3. If cost $> 0 \ \Rightarrow$ **LOP infeasible**; stop.
4. If cost $= 0$ and no artificial variable is in the basis, then a BFS was found.
5. Else, all $y_i^* = 0$, but some are still in the basis. Say we have $\boldsymbol{A}_{B(1)}, \ldots, \boldsymbol{A}_{B(k)}$ in basis $k < m$. There are $m - k$ additional columns of $\boldsymbol{A}$ to form a basis.

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6. Drive artificial variables out of the basis: If lth basic variable is artificial examine lth row of $\boldsymbol{B}^{-1}\boldsymbol{A}$. If all elements $= 0 \Rightarrow$ row redundant. Otherwise pivot with $\neq 0$ element.

2.2 Example

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$$\begin{array}{rrcrcrcrcl} \min & x_1 & + & x_2 & + & x_3 & & & & \\ s.t. & x_1 & + & 2x_2 & + & 3x_3 & & & = & 3 \\ & -x_1 & + & 2x_2 & + & 6x_3 & & & = & 2 \\ & & & 4x_2 & + & 9x_3 & & & = & 5 \\ & & & & & 3x_3 & + & x_4 & = & 1 \\ & \multicolumn{9}{l}{x_1, \ldots, x_4 \geq 0.} \end{array}$$

$$\begin{array}{lrcrcrcrcrcrcrcrcl} \min & & & & & & & & & x_5 & + & x_6 & + & x_7 & + & x_8 & & \\ s.t. & x_1 & + & 2x_2 & + & 3x_3 & & & + & x_5 & & & & & & & = & 3 \\ & -x_1 & + & 2x_2 & + & 6x_3 & & & & & + & x_6 & & & & & = & 2 \\ & & & 4x_2 & + & 9x_3 & & & & & & & + & x_7 & & & = & 5 \\ & & & & & 3x_3 & + & x_4 & & & & & & & + & x_8 & = & 1 \\ & \multicolumn{17}{l}{x_1, \ldots, x_8 \geq 0.} \end{array}$$

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		x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8
	−11	0	−8	−21	−1	0	0	0	0
$x_5 =$	3	1	2	3	0	1	0	0	0
$x_6 =$	2	−1	2	6	0	0	1	0	0
$x_7 =$	5	0	4	9	0	0	0	1	0
$x_8 =$	1	0	0	3	1*	0	0	0	1

		x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8
	−10	0	−8	−18	0	0	0	0	1
$x_5 =$	3	1	2	3	0	1	0	0	0
$x_6 =$	2	−1	2	6	0	0	1	0	0
$x_7 =$	5	0	4	9	0	0	0	1	0
$x_4 =$	1	0	0	3*	1	0	0	0	1

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		x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8
	−4	0	−8	0	6	0	0	0	7
$x_5 =$	2	1	2	0	−1	1	0	0	−1
$x_6 =$	0	−1	2*	0	−2	0	1	0	−2
$x_7 =$	2	0	4	0	−3	0	0	1	−3
$x_3 =$	1/3	0	0	1	1/3	0	0	0	1/3

		x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8
	-4	-4	0	0	-2	0	4	0	-1
$x_5 =$	2	2*	0	0	1	1	-1	0	1
$x_2 =$	0	$-1/2$	1	0	-1	0	$1/2$	0	-1
$x_7 =$	2	2	0	0	1	0	-2	1	1
$x_3 =$	$1/3$	0	0	1	$1/3$	0	0	0	$1/3$

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		x_1	x_2	x_3	x_4	x_5	x_6	x_7	x_8
	0	0	0	0	0	2	2	0	1
$x_1 =$	1	1	0	0	$1/2$	$1/2$	$-1/2$	0	$1/2$
$x_2 =$	$1/2$	0	1	0	$-3/4$	$1/4$	$1/4$	0	$-3/4$
$x_7 =$	0	0	0	0	0	-1	-1	1	0
$x_3 =$	$1/3$	0	0	1	$1/3$	0	0	0	$1/3$

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		x_1	x_2	x_3	x_4
	*	*	*	*	*
$x_1 =$	1	1	0	0	$1/2$
$x_2 =$	$1/2$	0	1	0	$-3/4$
$x_3 =$	$1/3$	0	0	1	$1/3$

3 A complete Algorithm for LO

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Phase I:

1. By multiplying some of the constraints by -1, change the problem so that $\boldsymbol{b} \geq \mathbf{0}$.
2. Introduce $y_1, \ldots, y_m$, if necessary, and apply the simplex method to min $\sum_{i=1}^{m} y_i$.
3. If cost> 0, original problem is infeasible; STOP.
4. If cost$= 0$, a feasible solution to the original problem has been found.
5. Drive artificial variables out of the basis, potentially eliminating redundant rows.

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Phase II:

1. Let the final basis and tableau obtained from Phase I be the initial basis and tableau for Phase II.
2. Compute the reduced costs of all variables for this initial basis, using the cost coefficients of the original problem.
3. Apply the simplex method to the original problem.

3.1 Possible outcomes

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1. Infeasible: Detected at Phase I.
2. $\boldsymbol{A}$ has linearly dependent rows: Detected at Phase I, eliminate redundant rows.
3. Unbounded (cost$= -\infty$): detected at Phase II.
4. Optimal solution: Terminate at Phase II in optimality check.

4 The big-M method

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$$\begin{array}{ll} \min & \sum_{j=1}^{n} c_j x_j + M \sum_{i=1}^{m} y_i \\ \text{s.t.} & \boldsymbol{Ax} + \boldsymbol{y} = \boldsymbol{b} \\ & \boldsymbol{x}, \boldsymbol{y} \geq \boldsymbol{0} \end{array}$$

5 The Column Geometry

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$$\begin{array}{lrcl} \min & \boldsymbol{c}'\boldsymbol{x} & & \\ \text{s.t.} & \boldsymbol{Ax} & = & \boldsymbol{b} \\ & \boldsymbol{e}'\boldsymbol{x} & = & 1 \\ & \boldsymbol{x} & \geq & \boldsymbol{0} \end{array}$$

$$x_1 \begin{bmatrix} \boldsymbol{A}_1 \\ c_1 \end{bmatrix} + x_2 \begin{bmatrix} \boldsymbol{A}_2 \\ c_2 \end{bmatrix} + \cdots + x_n \begin{bmatrix} \boldsymbol{A}_n \\ c_n \end{bmatrix} = \begin{bmatrix} \boldsymbol{b} \\ z \end{bmatrix}$$

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6 Computational efficiency

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Exceptional practical behavior: linear in n
Worst case

$$\begin{array}{ll} \max & x_n \\ \text{s.t.} & \epsilon \leq x_1 \leq 1 \\ & \epsilon x_{i-1} \leq x_i \leq 1 - \epsilon x_{i-1}, \qquad i = 2, \ldots, n \end{array}$$

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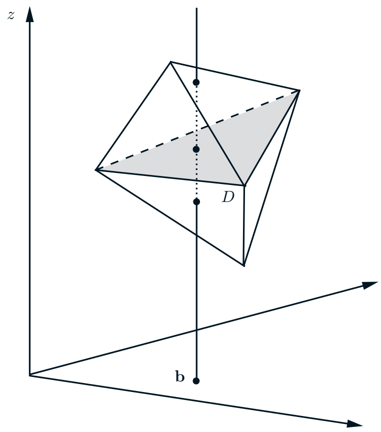
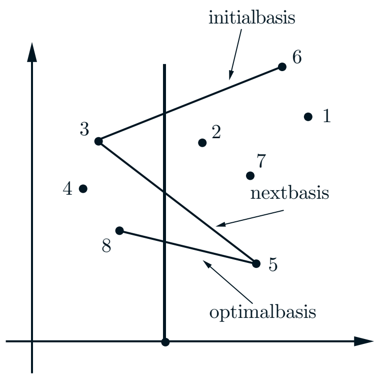
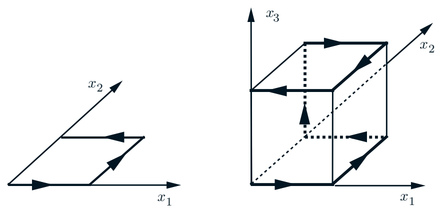
Theorem SLIDE 20


z
D
b


initialbasis
6
1
3
2
7
4
nextbasis
8
5
optimalbasis


x2
x1
x3
x2
x1

(a) (b)

- The feasible set has 2^n vertices
- The vertices can be ordered so that each one is adjacent to and has lower cost than the previous one.
- There exists a pivoting rule under which the simplex method requires $2^n - 1$ changes of basis before it terminates.

7 The Diameter of polyhedra

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- Given a polyhedron P, and $\boldsymbol{x}, \boldsymbol{y}$ vertices of P, the distance $d(\boldsymbol{x}, \boldsymbol{y})$ is the minimum number of jumps from one vertex to an adjacent one to reach $\boldsymbol{y}$ starting from $\boldsymbol{x}$.
- The diameter $D(P)$ is the maximum of $d(\boldsymbol{x}, \boldsymbol{y})$ $\forall \boldsymbol{x}, \boldsymbol{y}$.

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- $\Delta(n, m)$ as the maximum of $D(P)$ over all *bounded* polyhedra in $\Re^n$ that are represented in terms of m inequality constraints.
- $\Delta_u(n, m)$ is like $\Delta(n, m)$ but for possibly unbounded polyhedra.

7.1 The Hirsch Conjecture

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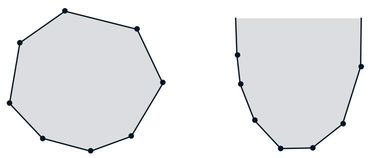
- $$\Delta(2, m) = \left\lfloor \frac{m}{2} \right\rfloor, \qquad \Delta_u(2, m) = m - 2$$



- **Hirsch Conjecture:** $\Delta(n, m) \leq m - n$.

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- We know that
$$\Delta_u(n, m) \geq m - n + \left\lfloor \frac{n}{5} \right\rfloor$$
$$\Delta(n, m) \leq \Delta_u(n, m) < m^{1+\log_2 n} = (2n)^{\log_2 m}$$

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