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6.006 Introduction to Algorithms
Spring 2008

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6.006 Recitation

Build 2008.25

6.006 Proudly Presents

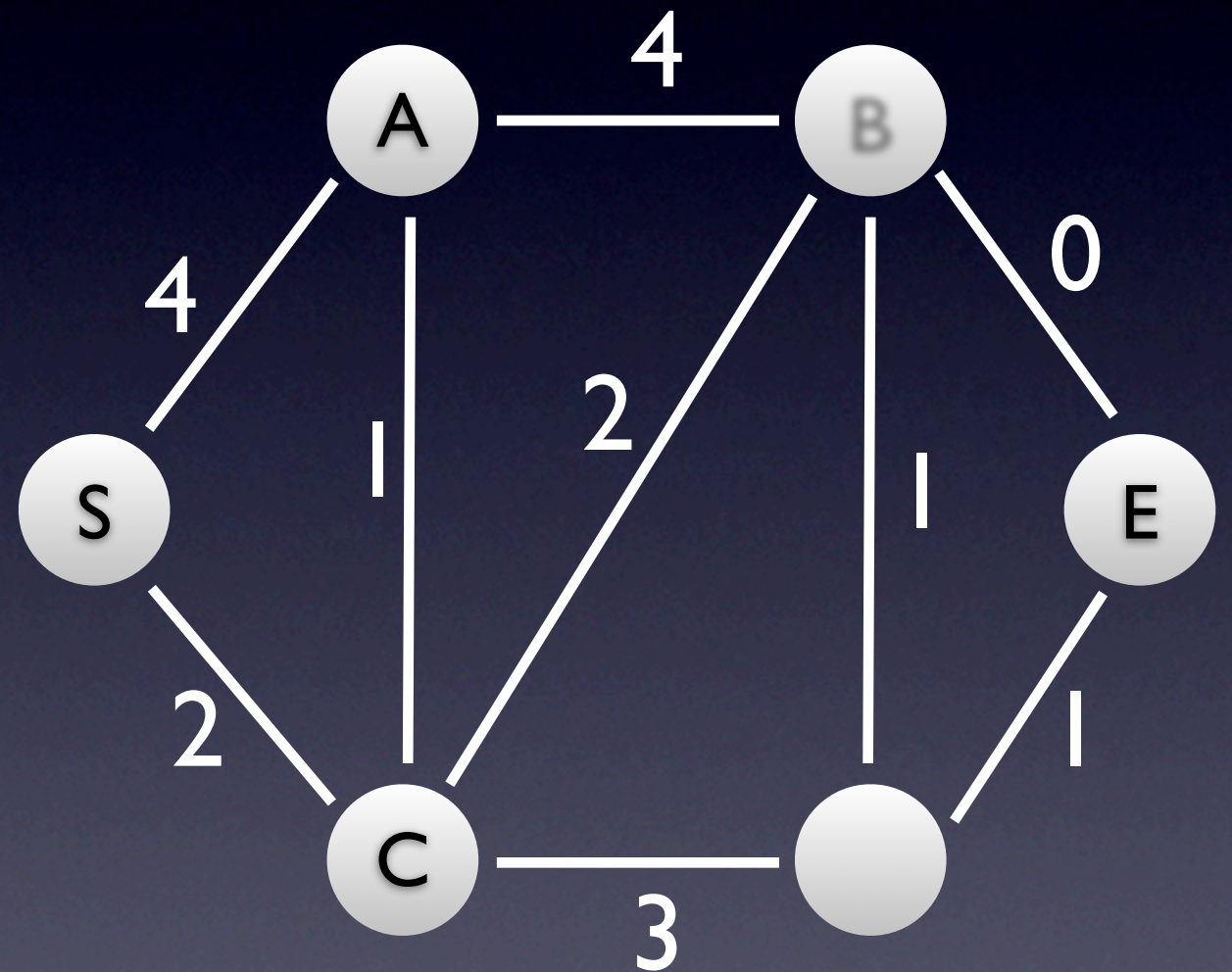
- Dijkstra: minimum-cost paths on crack
 - Algorithm
 - Concepts
 - Implementation
- Data structures come back from the dead (not talking about the quiz)

Minimum-Path Problem

- Given: graph G , source vertex s , edge costs
- Want: paths from s to everything else with minimum costs (sum of edge costs)
- Approach: let $d[v]$ be upper bounds for the real minimum costs, $\delta[v]$
- Start out easy: $d[v] = \infty$, $d[s] = 0$
- Relax until values in d converge to δ

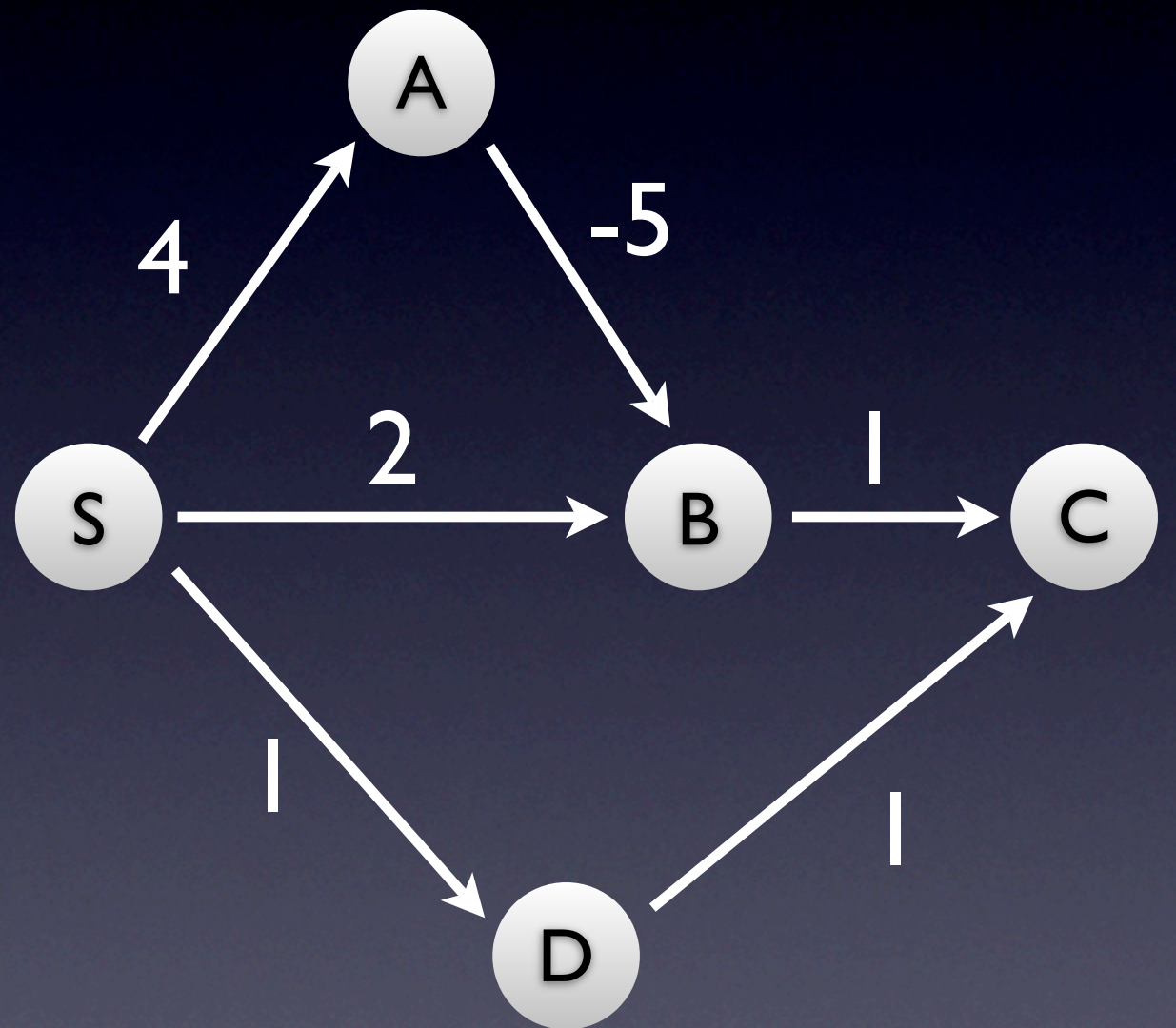
Good Dijkstra

- Generic initialization
- $U = V$
- Choose $v = \operatorname{argmin} d[v]$ in U , remove v from U
- Notice $d[v] = \delta[v]$
- Relax v 's outgoing edges
- Rinse, repeat



Bad Dijkstra

- Generic initialization
- $U = V$
- Choose $v = \operatorname{argmin} d[v \text{ in } U]$, remove v from U
- Notice $d[v] = \delta[v]$
- Relax v 's outgoing edges
- Rinse, repeat



Dijkstra Overview

- Nice and fast (that's why it's on crack)
- With limitations (crack impacts judgement)
 - Doesn't handle negative-cost edges
 - DOES handle 0-cost edges
 - Harder to code than Bellman-Ford

Dijkstra Works: Intuition

Dijkstra Works: Formal

Making Dijkstra Fast (its crack dealer)

- Generic initialization:
 $d[v] \leftarrow \infty, d[S] = 0$
- Choose $v = \operatorname{argmin} d[v]$,
by now $d[v] = \delta[v]$
- Relax all edges going out
of v
- Rinse, repeat
- Computing argmin
 - V times
- Relaxing
 - E times
- Looks like we need a
Data Structure

Min-Priority Queues

- Data Structure
 - **insert(key)** : adds to the queue
 - **min()** : returns the minimum key
 - **delete-min()** : deletes the min key
 - **delete(key)** : deletes the given key
 - optional (only needed in some apps)

Priority Queues with Min-Heaps

- Costs (see above line for explanations)
 - insert: $O(\log(N))$
 - min: $O(1)$
 - delete-min: $O(\log(N))$
 - delete: $O(\log(N))$ - only if given the index of the node containing the key

Priority Queues with PS3

- Is this priority queue monotone?
- Profit

Cool Python: Generators

1. Iterators

- used in for loops
- objects implementing `next()`

2. Generators

- express iterator functionality in a cooler way

```
1 def counter():
2     i = 0
3     while True:
4         yield i
5         i += 1

----
c = counter()
c.next()
>> 0
c.next()
>> 1
d = counter()
d.next()
>> 0
c.next()
>> 2
d.next()
>> 1
c.next()
>> 3
```


Dijkstra-Ready Priority Queues

```
1 class heap_id:
2     def __init__(self):
3         self.A = [None]
4         self.heapsize = 0
5         self.ID_to_index = {}
6         self._ID = self._ID_generator()
7     def insert(self, key):
8         """Returns an ID that is associated with the item."""
9         self.heapsize += 1
10        ID = self._ID.next()
11        self.ID_to_index[ID] = self.heapsize
12        self.A.append( [positive_infinity(), ID] )
13        self.decrease_key(self.heapsize, key)
14        return ID
15    def _ID_generator(self):
16        ID = 0
17        while True:
18            yield ID
19            ID += 1
```


Dijkstra-Ready Priority Queues II

```
1 class heap_id:
2     def decrease_key_using_id(self, ID, key):
3         """Decrease key given ID."""
4         self.decrease_key(self.ID_to_index[ID], key)
5     def extract_min(self):
6         """Extracts min and returns key."""
7         return self.extract_min_with_id()[0]
8     def extract_min_with_id(self):
9         """Extracts min and returns (key,ID) pair."""
10        if self.heapsize < 1:
11            print "error: empty heap"
12            return
13        self._swap(1, self.heapsize)
14        self.heapsize -= 1
15        min_pair = self.A.pop()
16        del self.ID_to_index[min_pair[1]]
17        self.min_heapify(1)
18        return tuple(min_pair)
```